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2 METHOD AND SYSTEM FOR TRANSMITTING DATA BETWEEN A
3 RECEIVER AND A TRANSMITTER

4 The present invention generally relates to a method and a system
5 for transmitting data between at least one receiver operatively connected to at
6 least one transmitter, and more particularly to a method and a system for
7 permissible transmission via at least one high-speed link having a plurality of
8 virtual channels.

9 Virtual channels are generally used for maximum efficiency
10 usage of a single bi-directional serial link between a receiver and a transmitter.
11 The communication is usually based on asynchronous concurrent processes at
12 the packet level. More specifically, Communicating Sequential Processes
13 ("CSP") is used for synchronized communication between the receiver and
14 transmitter, meaning each communicating process waits until data transfer is
15 complete before continuing.

1 One known method for maintaining the CSP while ensuring that
2 no data is lost or overwritten suggests the use of acknowledgement packets.
3 The receiver must acknowledge each packet before another can be sent, and all
4 the packets of a message must also be acknowledged before the sending
5 process continues. In this case, when the data packet is received by a process,
6 an acknowledgement packet is sent back to the transmitter. The
7 acknowledgement packet allows the transmitter to send another packet after a
8 packet has been received or even partially received.

9 The problem with using acknowledgement packets is that the
10 transmitter does not necessarily know if the receiver is ready to receive a
11 packet in the first place. When the transmitter sends a packet to the receiver
12 that does not have any available buffer, the packet is either lost or the link is
13 flow controlled at the physical layer. If the packet is lost, the transmitter will
14 never receive the acknowledgement packet. As a result, it must rely on some
15 timer to know when to try again, causing delay and waste of bandwidth. On
16 the other hand, if the data is flow controlled at the physical layer, data
17 transmission is indefinitely postponed until the receiver is able to accept the
18 data. Again, this causes both delay and waste of bandwidth by blocking other
19 virtual channels.

20 Accordingly, it is an object of the present invention to provide an
21 improved method and system for transmitting data using a virtual channel
22 credit packet over a high-speed link with a plurality of virtual channels.

23 A related object of the present invention is to provide an
24 improved method and system for transmitting data between a receiver and
25 transmitter without the use of acknowledgement packets.

1 Another object of the present invention is to provide an improved
2 method and system for allowing a transmitter to send a data packet to a receiver
3 only after the receiver has resources to receive the packet.

4 Still another object of the present invention is to provide an
5 improved method and system for transmitting data via at least one high-speed
6 link with a plurality of virtual channels using more efficient bandwidth.

7 A further object of the present invention is to provide an
8 improved method and system for transmitting data using a virtual channel
9 credit packet having a unique assigned virtual channel number for each virtual
10 channel.

11 BRIEF SUMMARY OF THE INVENTION

12 The present invention provides a method and a system for
13 transmitting data between one or more receivers operatively connected to one
14 or more transmitters, and more particularly to a method and a system for
15 permissible transmission via at least one high-speed link having a plurality of
16 virtual channels. By having the receiver send a credit packet with a unique
17 virtual channel number that is specifically assigned to each virtual channel only
18 when it has a buffer available, the present invention is able to bypass the need
19 for an acknowledgement packet, resulting in reliable transmissions and
20 efficient use of bandwidth.

21 In accordance with this invention, the receiver first sends a virtual
22 channel credit packet for a particular virtual channel to the transmitter only if
23 the receiver is available to receive data. A unique virtual channel number is
24 assigned for that particular virtual channel, and is included in the virtual
25 channel credit packet. Then, the transmitter responds to the virtual channel

1 credit packet accordingly, which includes transmitting data to the receiver if it
2 is actually available. The receiver then receives the data packet transmitted
3 from the transmitter. The sending of a credit packet guarantees that the
4 receiver is ready and has the available resources to receive transmission data
5 from the transmitter.

6 Other objects, features and advantages will become apparent
7 upon reading the following detailed description, in conjunction with the
8 attached drawings, in which:

9 FIGURE 1 is a schematic diagram of an exemplary connection
10 between a peripheral and an I/O card in which the present invention may be
11 implemented;

12 FIG. 2 is an ensemble illustrating the format and control
13 characters for the data packet and the credit packet for one virtual channel;

14 FIG. 3 is a sequence illustrating the flow of the credit packet and
15 the data packet for one virtual channel;

16 FIG. 4 is a flowchart illustrating the subroutine executed by the
17 receiver; and,

18 FIG. 5 is a flowchart illustrating the subroutine executed by the
19 transmitter.

20 DETAILED DESCRIPTION

21 Broadly stated, the present invention is directed to a method and a
22 system for transmitting data between at least one receiver operatively
23 connected to at least one transmitter over a high-speed link with a plurality of
24 virtual channels. Each virtual channel is assigned with a unique virtual channel
25 number. When the receiver is ready for transmission for a particular virtual

1 channel, it sends a virtual channel credit packet bearing the assigned virtual
2 channel number. The transmitter then responds to the virtual channel credit
3 packet. After the transmission for this particular virtual channel is finished, the
4 process is repeated for the next virtual channel until all the virtual channels are
5 running.

6 Turning now to FIG. 1, a schematic diagram of an exemplary
7 connection between a peripheral device or peripheral and an I/O card is shown
8 and illustrates one way in which the receiver is connected to the transmitter for
9 the implementation of the present invention. However, it should be understood
10 that other connections are possible and are within the scope of the present
11 invention. In this example, a peripheral 10 is connected to an I/O card 12
12 through a bi-directional serial link 14 with a plurality of virtual channels 16. In
13 this case, the peripheral 10 is the transmitter, and the I/O card 12 is the
14 receiver. Although any number of virtual channels 16 can be used, 256 virtual
15 channels are preferred for this implementation. The virtual channels can start
16 from the number zero, and end with any desired number N as long as it is
17 finite. The only limitation is the hardware. However, the number of virtual
18 channels is important in that it determines the number of credit packets that
19 will be used in the present invention, as will be described. Furthermore, the
20 present method loops to repeat the process until all the virtual channels are
21 running.

22 The preferred ensemble of a virtual channel data packet indicated
23 generally at 20 and a credit packet indicated generally at 22 for one virtual
24 channel is shown in detail in FIGs. 2 and 3. Each virtual channel is assigned a
25 unique Virtual Channel Number (VCN) 24. As shown in FIG. 2, the packets
26 include some of the control characters 26 defined by the Institute of Electrical
27 and Electrics Engineers Standards (IEEE Std) 1355-1995 for Heterogeneous

1 Interconnect. In addition, the VCN 24 is used for both the data packet and the
2 credit packet formats for each virtual channel. Because there is a number
3 assigned for each virtual channel, the transmitter 10 and receiver 12 are able to
4 keep track of the transmission within the plurality of virtual channels 16.

5 When the receiver 12 sends a virtual channel credit packet 22 for
6 a particular virtual channel, the transmitter 10 knows to which virtual channel
7 the credit packet is designated. In other words, the transmitter 10 is now
8 granted permission to send a data packet 20, if available, to the receiver 12
9 through the virtual channel designated by the credit packet 22. In FIG. 3, the
10 sequence of the flow of the credit packet 22 and the data packet 20 for a single
11 virtual channel is illustrated, with the direction of the arrows beneath the packet
12 information depicting the direction of packet flow between transmitter and
13 receiver. Each data packet 20 from a transmitter can be transmitted only when
14 a credit packet 22 is issued by the receiver 12. FIG. 3 shows the basic flow for
15 multiple transmissions that may take place for each virtual channel.

16 In accordance with an important aspect of the present invention, a
17 flowchart illustrating the subroutine of the receiver 12 in FIG. 4. In order for
18 the receiver 12 to initially start the process (block 28) for a plurality of virtual
19 channels 16, there must be an operable communication link established (block
20 30) between the transmitter 10 and the receiver. Once that is done, the receiver
21 12 spawns a receiver process for a particular virtual channel number N (block
22 32). The receiver 12 first checks if there is any available buffer for VCN "N"
23 (block 34). If the receiver 12 does not have any buffer available for
24 transmission (block 36), it waits for a predetermined time and rechecks for
25 available buffer until it finds some buffer (block 34). Once available buffer is
26 found (block 36), then the receiver 12 sends a virtual channel credit packet 22
27 for VCN "N" to the transmitter (block 38).

1 In contrast, if there is available buffer (block 54), the transmitter
2 10 next looks for the virtual channel credit packet 22 with its assigned VCN
3 "N" from the receiver 12 (block 56). If a credit packet 22 can not be found, the
4 transmitter waits for a predetermined time and rechecks for the credit packet 22
5 for that particular virtual channel with VCN "N" (block 56). If there is a credit
6 packet 22, it checks to determine if it has any data packets 20 for transmission
7 (block 58). The transmitter again waits for a predetermined time and keeps
8 checking for data packets 20 if there is no data packet for transmission (block
9 58). However, if there is a data packet 20, with the VCN "N" as described
10 earlier, it will accordingly be transmitted to the receiver 12 (block 60). After
11 the transmission of the data packet 20 for VCN "N," the transmitter 10 loops
12 back to continue checking for an available buffer for VCN "N" (block 52).
13 However, it also repeat the transmitter process for the next VCN or VCN
14 "N+1" until all the VCNs are running (block 62). When all the VCNs are
15 running, the transmitter process will exit the subroutine (block 64).

